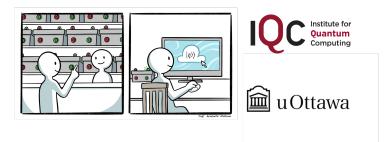
Quantum delegation with an off-the-shelf device

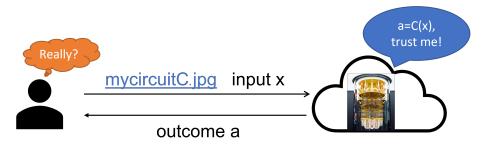
Arthur Mehta

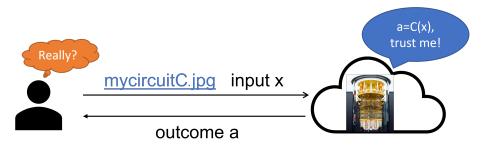
University of Ottawa

Joint work with Anne Broadbent and Yuming Zhao, based on arXiv:2304.03448



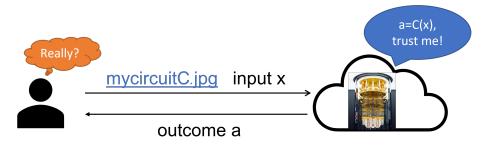






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- I want to be convinced of the correctness,
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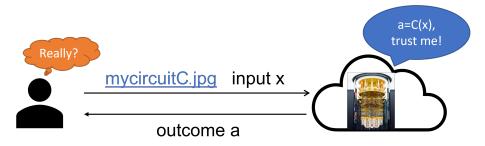


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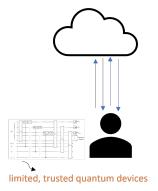


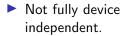
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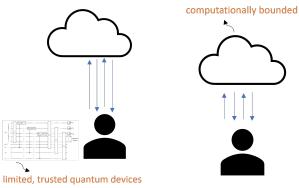
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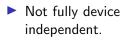
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- Goal: zero-knowledge verifiable delegation of quantum computations

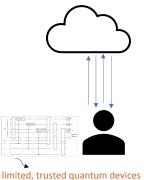








 Computational and hardness assumptions.

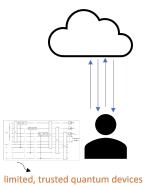


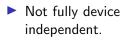


 Not fully device independent. Computational and hardness assumptions.

computationally bounded

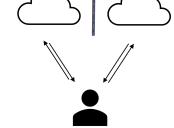
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Our model: a single quantum server + an untrusted device, all in a single round

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Set-up: Purchase an off-the-shelf based on |x|.

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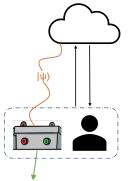
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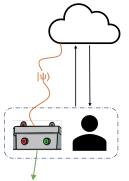
Note The shared state only depends on |x|





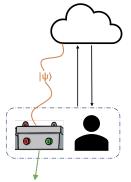


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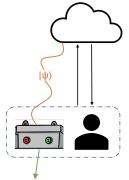


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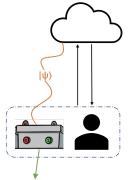


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A measurement device

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- 1. Circuit-to-Hamiltoniam construction $C \rightarrow H: C$ accepts $\Leftrightarrow \lambda_0(H)$ is small
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3. Estimates the ground energy of H



- Accepts if $\lambda_0(H)$ is low
- Rejects if $\lambda_0(H)$ is high

Main Results

Theorem (Broadbent, M, Zhao 2023)

All efficient quantum computations have verifiable delegation protocols in OTS.

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Can amplified to constant completeness-soundness gap while preserving ZK.





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Verifier's Task: Certify many EPR pairs.



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Theorem

Very Informal: Prove an enhanced version of Gowers Hatami theorem from approximate representation theory.

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 - The LWPBT has alot of entanglement left over after the test.
- 2. We know $QMA \subseteq OTS$. Is it possible $OTS = MIP^*$?
 - What if we lift the constant measurement requirement?
- 3. Applications to PoQK via self-testing.
 - Can we show our overall protocol self-tests for ground states?